A Kochen-Specker system has at least 21 vertices

Sander Uijlen suijlen@cs.ru.nl

Bas Westerbaan bwesterb@cs.ru.nl

Radboud Universiteit Nijmegen

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A Kochen-Specker system S is a finite set of points on the (open) northern hemisphere, such that there is no 010-coloring; that is: there is no $\{0, 1\}$ -valued coloring with

- 1. three pairwise orthogonal points are assigned (1,0,0), (0,1,0) or (0,0,1) and
- 2. two orthogonal points are not assigned (1, 1).

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Theorem (Kochen-Specker)

There is a Kochen-Specker system. Thus: there is no non-contextual deterministic theory predicting the measurement of a SPIN-1 particle.

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Penrose, Peres (indep.)	1991	
Conway	~ 1995	

Arends, Wampler, Ouaknine 2009

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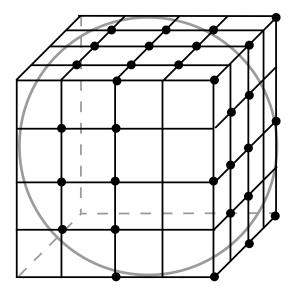
Arends, Wampler, Ouaknine $2009 \ge 18$

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U&W		may	≥ 21
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Conway	~ 1995	\leq 31
U&W	july?	$\geq 22 \text{ or} = 21$
U&W	may	≥ 21
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Conway's record



It is a problem about graphs

Given a finite set of points S on the projective plane, its orthogonality graph $\mathcal{G}(S)$ has as vertices the points and two points are adjacent if and only if they are orthogonal. Given a finite set of points S on the projective plane, its orthogonality graph $\mathcal{G}(S)$ has as vertices the points and two points are adjacent if and only if they are orthogonal.

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A graph G is embeddable if there is a S such that $G \leq \mathcal{G}(S)$.

A 010-coloring of a graph, is a $\{0,1\}$ -vertex coloring, such that

- 1. every triangle is colored (1,0,0), (0,1,0) or (0,0,1) and
- 2. no adjacent vertices are colored both 1.

It is a problem about graphs

There is a Kochen-Specker system with n points if and only if there is a embeddable and non-010-colorable graph on n vertices.

(The orthogonality graph of) a minimal Kochen-Specker system is connected; $\sim 10^{26.4}$

Restrict the search

(The orthogonality graph of) a minimal Kochen-Specker system is connected; $\sim 10^{26.4}$ v squarefree and $\sim 10^{10.2}$

The candidates

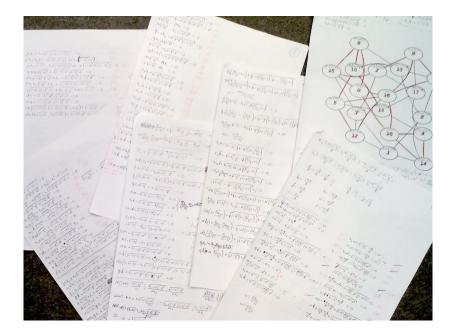
Our computation found the following number of non 010-colorable squarefree graphs with minimal vertex degree 3

#V	# candidates
\leq 16	0
17	1
18	2
19	19
20	441

The candidates

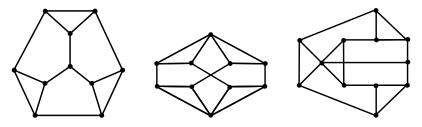
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20	441
21	\geq 7616

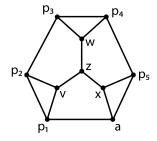


Unembeddable subgraphs

All these candidates contain as a subgraph one of these unembeddable graphs.



Suppose this graph is embeddable.



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 p_2 p_1 p_4 Note distin p_2 p_5 Simila And s

Note that v and a are distinct points orthogonal to p_1 . Thus p_1 is fixed. Observe: p_1 is collinear to $v \times a$.

Similarly: p_2 is collinear to $v \times (v \times a)$. And so on. We see *a* is collinear to $x \times (x \times (w \times (w \times (v \times (v \times a)))))$.

We may assume z = (0, 0, 1), x = (1, 0, 0), $v = (v_1, v_2, 0)$, $w = (w_1, w_2, 0)$ and $a = (0, a_2, a_3)$. We have:

$$\begin{pmatrix} 0\\ a_2\\ a_3 \end{pmatrix} \text{ is collinear to} \begin{pmatrix} 0\\ -a_2v_1w_2(v_1w_1+v_2w_2)\\ -a_3(v_1^2w_1^2+v_1^2w_2^2+v_2^2w_1^2+v_2^2w_2^2) \end{pmatrix}$$

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That is:

$$\begin{split} v_1 w_2 \langle v, w \rangle &= v_1 w_2 (v_1 w_1 + v_2 w_2) \\ &= v_1^2 w_1^2 + v_1^2 w_2^2 + v_2^2 w_1^2 + v_2^2 w_2^2 \\ &= (v_1^2 + v_2^2) w_1^2 + (v_1^2 + v_2^2) w_2^2 \\ &= w_1^2 + w_2^2 = 1. \end{split}$$

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Since v and w are not collinear, we have by Cauchy-Schwarz $|\langle v, w \rangle| < 1$. Note $|v_1|, |w_2| \le 1$. Thus: $|v_1w_2 \langle v, w \rangle| < 1$. Contradiction.

Example of automized cross product chasing

```
load_package redlog;
rlset R;
procedure d(x.v);
    (first x) * (first y) +
    (second x) * (second v) +
    (third x) * (third v);
procedure k(x.v);
    {(second x)*(third y) - (third x)*(second y),
     (third x)*(first v) - (first x)*(third v).
     (first x)*(second v) - (second x)*(first v)};
v0c1 := 1; v0c2 := 0; v0c3 := 0;
v1c1 := 0; v1c2 := 1; v1c3 := 0;
v0 := {v0c1, v0c2, v0c3}:
v1 := {v1c1, v1c2, v1c3}:
v2 := \{v2c1, v2c2, v2c3\};
v3 := {v3c1, v3c2, v3c3};
v2c1 := 0:
neq0 := k(v0,k(v3,v1));
```

```
(snip)
```

```
(snip)
```

```
(first neq29 neq 0 or
second neq29 neq 0 or
third neq29 neq 0 or
d(v2,v0) = 0 and
d(k(v3,v0),v3),k(k(k(k(v3,v1),v1),v2),v2)) = 0 and
true;
rlqe ex(v3c3,
ex(v3c4,
ex(v3c4,
ex(v3c4,
ex(v2c2,phi)))));
```

Source code, paper and experimental results can be found at

kochen-specker.info

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```

Some open problems:

- If G is embeddable, is there a S such that $G = \mathcal{G}(S)$.
- Is every embeddable graph, grid embeddable? That is: using points of the form (^x/_{√n}, ^y/_{√n}, ^z/_{√n}) for x, y, z, n ∈ Z.